Math 2603 - Lecture 19 Section 9.3 Isomorphism of Graphs

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Motivation

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Different visualization of the same graph

Remark

As we explained last time, the essential feature of a graph is the incidental relationship among its vertices. Therefore, there are different visualization of the same graph.

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Example: Petersen graph



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What about these two graphs? They are both the Petersen graph.

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Remark

We need a definition of two graphs being essential the same.

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Remark

In mathematics we have several terms to address different kinds of "equality". Here we mean two graphs that are essentially the same. In other words, they have the same structures. The term we use is **isomoprhic**.

Graph Isomorphism

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Definition

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Given graphs $\mathcal{G}_1 = (\mathcal{V}_1, \mathcal{E}_1)$ and $\mathcal{G}_2 = (\mathcal{V}_2, \mathcal{E}_2)$, we say that \mathcal{G}_1 is isomorphic to \mathcal{G}_2 and write $\mathcal{G}_1 \cong \mathcal{G}_2$, if and only if there is a one-to-one function $\varphi : \mathcal{V}_1 \to \mathcal{V}_2$ that induces a one-to-one function $\mathcal{E}_1 \to \mathcal{E}_2$ as well. In other words,

• if vw is an edge in \mathcal{E}_1 , then $\varphi(v)\varphi(w)$ is also an edge in \mathcal{E}_2 ;

• every edge in \mathcal{E}_2 is of the form $\varphi(v)\varphi(w)$ where $vw \in \mathcal{E}_1$. We call φ an **isomorphism** from \mathcal{G}_1 to \mathcal{G}_2 . Abusing of notation, we also say $\varphi: \mathcal{G}_1 \to \mathcal{G}_2$ is an isomorphism.

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Graph isomorphism is an equivalence relation

Proposition

On the set of all graphs, the relation \cong is an equivalence relation.

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Sketch of proof.

Reflexivity: the identity map works. Symmetry: if $\varphi : \mathcal{G}_1 \to \mathcal{G}_2$ is an isomorphism, so is $\varphi^{-1} : \mathcal{G}_2 \to \mathcal{G}_1$. Transitivity: if $\varphi_1 : \mathcal{G}_1 \to \mathcal{G}_2$ and $\varphi_2 : \mathcal{G}_2 \to \mathcal{G}_3$ are both isomorphisms, then $\varphi_2 \circ \varphi_1$ is an isomorphism from \mathcal{G}_1 to \mathcal{G}_3 . \Box

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Necessary conditions for graph isomorphism

Proposition

If \mathcal{G}_1 and \mathcal{G}_2 are isomorphic graphs, then they have the

- same number of vertices;
- same number of edge;
- same degree sequences.

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Remark

This proposition is very useful to identify non-isomorphic graphs.

Example: which graphs are isomorphic?



Example: which graphs are isomorphic

Solution

Graph	V	E	Deg	~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~	\sim	٨٩	
(i)	4	5	3, 3, 2, 2		\sim		
<i>(ii)</i>	4	4	2, 2, 2, 2	(i)	(ii)	(iii)	
(iii)	4	5	3, 3, 2, 2	a p	$\langle \rangle$		
(iv)	4	4	3, 2, 2, 1	\checkmark	\checkmark		
(v)	4	5	3, 3, 2, 2	(iv)	(v)		

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Solution

Graph	V	Ε	Deg]	~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~	<u>م</u>	۶
(i)	4	5	3, 3, 2, 2			\sim	
(ii)	4	4	2, 2, 2, 2]	(i)	(ii)	(iii)
(iii)	4	5	3, 3, 2, 2		a p	÷	
(iv)	4	4	3, 2, 2, 1		\checkmark	\checkmark	
(v)	4	5	3, 3, 2, 2		(iv)	(v)	

It turns out that (i), (iii), (v) are all isomorphic, and (ii) and (iv) are not isomorphic to any other one.

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More properties preserved by graph isomorphisms

Remark

There are a lot more properties of graphs that are preserved by graph isomorphisms. As a result, if two graphs have different status with respect to some property, then they are not isomorphic.

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There are a lot more properties of graphs that are preserved by graph isomorphisms. As a result, if two graphs have different status with respect to some property, then they are not isomorphic.

- connectivity;
- length of the longest path/cycle;
- whether it has an Eulerian path/circuit;
- whether it has a Hamiltonian path/cycle;
- whether it is planar.

We will elaborate them in the next few weeks.

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Homework Assignment #11

Section 9.3 Exercise 3(a), 5(b), 8, 11(a).

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